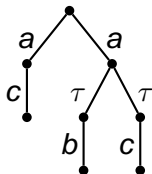
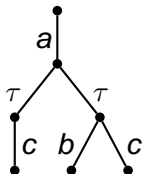
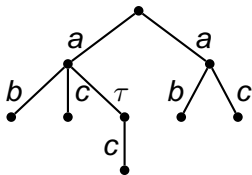
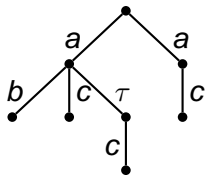
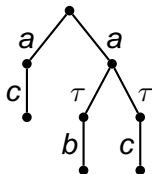
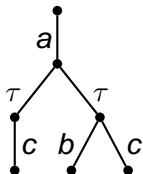
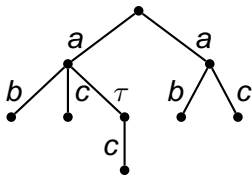
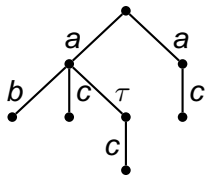


## Good old days: Observing trees in JCMB



For fun: add probabilities

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For fun: add probabilities

# Some remarks on testing probabilistic processes

Matthew Hennessy

and Yuxin Deng, Rob van Glabbeek, Carroll Morgan, Chenyi Zhang  
*Nicta*

gdp festschrift

# Outline

## Testing theory

Probabilistic CSP

The power of testing

Simulations

Equational theories

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## Testing theory

### Probabilistic CSP

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## Testing scenario

- ▶ a set of processes  $\mathcal{Proc}$
- ▶ a set of tests  $\mathcal{T}$
- ▶ a set of outcomes  $\mathcal{O}$
- ▶  $Apply : \mathcal{T} \times \mathcal{Proc} \rightarrow \mathcal{P}_{fin}^+(\mathcal{O})$  – the non-empty finite set of possible results of applying a test to a process

### Comparing sets of outcomes:

- ▶  $O_1 \sqsubseteq_{H_0} O_2$  if for every  $o_1 \in O_1$  there exists some  $o_2 \in O_2$  such that  $o_1 \leq o_2$
- ▶  $O_1 \sqsubseteq_{S_m} O_2$  if for every  $o_2 \in O_2$  there exists some  $o_1 \in O_1$  such that  $o_1 \leq o_2$

$o_1 \leq o_2$  : means  $o_2$  is as least as good as  $o_1$

## Testing scenario

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## Testing preorders

- ▶  $P \sqsubseteq_{\text{may}} Q$  if  $\text{Apply}(T, P) \sqsubseteq_{\text{Ho}} \text{Apply}(T, Q)$  for every test  $T$
- ▶  $P \sqsubseteq_{\text{must}} Q$  if  $\text{Apply}(T, P) \sqsubseteq_{\text{Sm}} \text{Apply}(T, Q)$  for every test  $T$

### Standard testing:

Use as outcomes  $\mathcal{O} = \{\top, \perp\}$  with  $\perp \leq \top$

### Probabilistic testing:

Use as  $\mathcal{O}$  the unit interval  $[0, 1]$

Intuition: with  $0 \leq p \leq q \leq 1$ , passing a test with probability  $q$  better than passing with probability  $p$

### Theorem

For  $O_1, O_2 \in \mathcal{P}_{\text{fin}}^+(\mathcal{O}_{\text{prob}})$  we have

- ▶  $O_1 \sqsubseteq_{\text{Ho}} O_2$  if and only if  $\max(O_1) \leq \max(O_2)$
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# Finite probabilistic CSP: pCSP

- ▶ **0** Do nothing
- ▶  $a.P$  Perform  $a$  then act as  $P$
- ▶  $P \mid_A Q$  Run  $P$  and  $Q$  in parallel ...
- ▶  $P \square Q$  External nondeterministic choice between  $P$  and  $Q$  - environment chooses
- ▶  $P \sqcap Q$  Internal nondeterministic choice between  $P$  and  $Q$
- ▶  $P_p \oplus Q$  Probabilistic choice: act like  $P$  with probability  $p$ , as  $Q$  with probability  $(1 - p)$

CSP: the sublanguage without probabilistic choice

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CSP: the sublanguage without probabilistic choice

# Operational semantics

## LTS:

a triple  $\langle S, \text{Act}_\tau, \rightarrow \rangle$ , with

- ▶  $S$  a set of states
- ▶  $\text{Act}_\tau$  a set of actions  $\text{Act}$ , with extra  $\tau$
- ▶  $\rightarrow \subseteq S \times \text{Act}_\tau \times S$  - the effect of performing actions.

## pLTS, probabilistic LTS:

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## (finite) Distributions $\mathcal{D}(S)$ :

Mappings  $\Delta : S \rightarrow [0, 1]$  with

$$\sum \text{DGHMZ (Nicta)} = 1$$

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# Operational semantics

## CSP as an LTS:

- ▶ States  $S$  are all terms from CSP
- ▶ Actions  $s \xrightarrow{\alpha} s'$  defined inductively

## pCSP as a pLTS:

- ▶ States  $S$  are **subset**  $S_p$  of terms from pCSP
- ▶ Actions  $s \xrightarrow{\alpha} s'$  defined inductively
- ▶ Terms  $P$  in pCSP interpreted as **distributions**  $\llbracket P \rrbracket$  over  $S$



# pCSP as a pLTS

## States $S_p$ :

- ▶  $\mathbf{0} \in S_p$
- ▶  $a.P \in S_p$
- ▶  $P \sqcap Q \in S_p$
- ▶  $s_1, s_2 \in S_p$  implies  $s_1 \square s_2 \in S_p$
- ▶  $s_1, s_2 \in S_p$  implies  $s_1 \mid_A s_2 \in S_p$ .

## Distributions $\llbracket P \rrbracket$ :

- ▶  $\llbracket \mathbf{s} \rrbracket = \bar{\mathbf{s}}$  one point distribution  $s \rightarrow 1$
- ▶  $\llbracket P_p \oplus Q \rrbracket = p \cdot \llbracket P \rrbracket + (1-p) \cdot \llbracket Q \rrbracket$
- ▶  $\llbracket P \square Q \rrbracket = \llbracket P \rrbracket \square \llbracket Q \rrbracket$
- ▶  $\llbracket P \mid_A Q \rrbracket = \llbracket P \rrbracket \mid_A \llbracket Q \rrbracket$

# pCSP as a probabilistic LTS

Defining arrows:  $S_p \xrightarrow{\alpha} \mathcal{D}(S_p)$

(ACTION)

$$a.P \xrightarrow{a} [P]$$

(EXT.I.L)

$$S_1 \xrightarrow{\tau} \Delta$$

$$\frac{S_1 \xrightarrow{\tau} \Delta}{S_1 \square S_2 \xrightarrow{\tau} \Delta \square \overline{S_2}}$$

(PAR.L)

$$S_1 \xrightarrow{\alpha} \Delta$$

$$\frac{S_1 \xrightarrow{\alpha} \Delta}{S_1 \mid_A S_2 \xrightarrow{\alpha} \Delta \mid_A \overline{S_2}} \quad \alpha \notin A$$

(INT.L)

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(PAR.I)

$$S_1 \xrightarrow{a} \Delta_1, S_2 \xrightarrow{a} \Delta_2$$

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and symmetric rules

Very similar to standard rules

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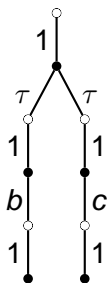
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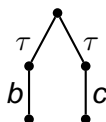
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# Example pLTSs



abbrev. to



$$b. \mathbf{0} \sqcap c. \mathbf{0}$$



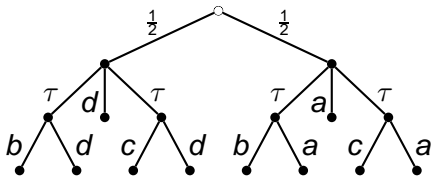
states



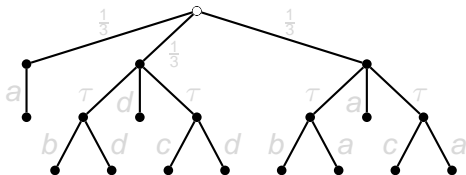
distributions

# Example pLTSs

$$(b \sqcap c) \sqcap (d \frac{1}{2} \oplus a)$$

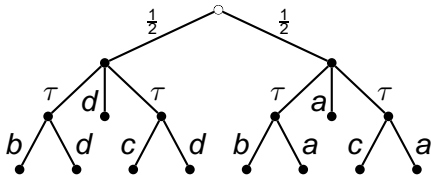


$$a \frac{1}{3} \oplus ((b \sqcap c) \sqcap (d \frac{1}{2} \oplus a))$$

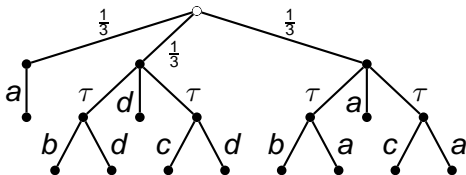


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$$(b \sqcap c) \sqcap (d \frac{1}{2} \oplus a)$$



$$a \frac{1}{3} \oplus ((b \sqcap c) \sqcap (d \frac{1}{2} \oplus a))$$



# Testing pCSP processes

## Tests:

Any process which may contain new *report success* action  $\omega$

$a.\omega \frac{1}{4} \oplus (b \square c.\omega)$ :

- ▶ 25% of time requests an  $a$  action
- ▶ 75% requests a  $c$  action
- ▶ 75% requires that  $b$  is not possible in a must test

## Applying test $T$ to process $P$ :

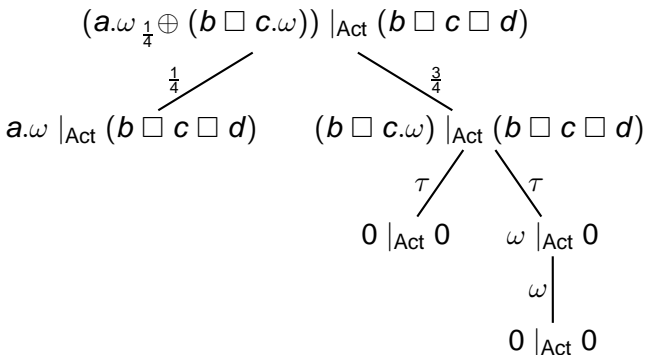
- ▶ Run the combined process  $T \mid_{\text{Act}} P$
- ▶ Calculate chances of success

Note:  $T \mid_{\text{Act}} P$  can only perform  $\tau$  or *report success*

## Running a test

$$\text{Test: } T = a.\omega \frac{1}{4} \oplus (b \square c.\omega)$$

$$\text{Process } P = b \square c \square d$$

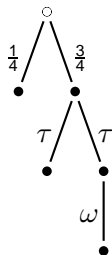




## Calculating chances of success

$$\text{Test: } T = a.\omega_{\frac{1}{4}} \oplus (b \square c.\omega)$$

$$\text{Process } P = b \square c \square d$$



Calculate success probabilities:

$$\text{Apply}(T, P) = \frac{1}{4} \cdot \{0\} + \frac{3}{4} \cdot \{0, 1\} = \{0, \frac{3}{4}\}$$

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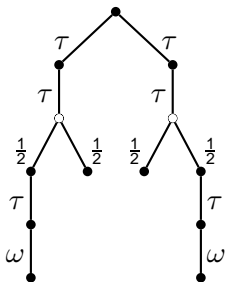
**The power of testing**

Simulations

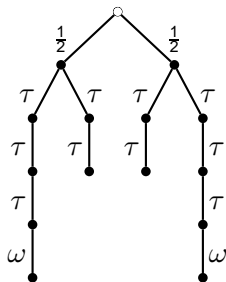
Equational theories

## Example: probabilistic choice and prefixing

$$R_1 = a.(b \frac{1}{2} \oplus c) \quad R_2 = a.b \frac{1}{2} \oplus a.c \quad T = a.b.\omega \sqcap a.c.\omega$$



$$\text{Apply}(T, R_1) = \left\{ \frac{1}{2} \right\}$$

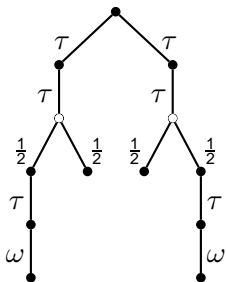


$$\text{Apply}(T, R_2) = \{0, \frac{1}{2}, 1\} = \frac{1}{2} \cdot \{0, 1\} + \frac{1}{2} \cdot \{0, 1\}$$

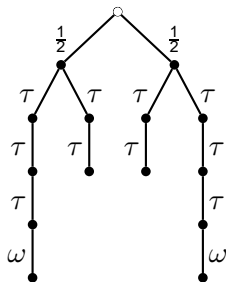
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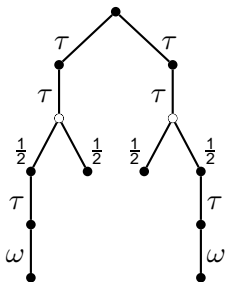


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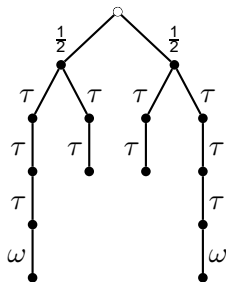
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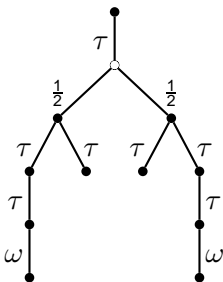
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## A beloved testing axiom

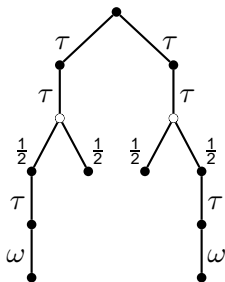
$$R_1 = a.(b \sqcap c)$$

$$R_2 = a.b \sqcap a.c$$

$$T = a.(b.\omega \frac{1}{2} \oplus c.\omega)$$



$$\text{Apply}(T, R_1) = \{0, \frac{1}{2}, 1\}$$



$$\text{Apply}(T, R_2) = \{\frac{1}{2}\}$$

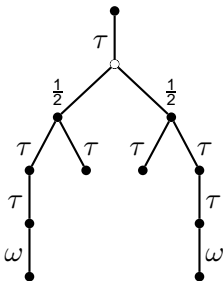
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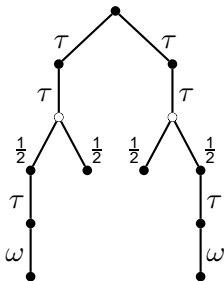
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So action prefix does **not** distribute over internal choice

## Power of internal choice in tests

- ▶  $R_1 = a \cdot \frac{1}{2} \oplus (b \square c)$        $R_2 = (a \square b) \cdot \frac{1}{2} \oplus (a \square c)$
- ▶  $T = a \cdot (\omega \cdot \frac{1}{2} \oplus \mathbf{0}) \sqcap (b \cdot \omega \cdot \frac{1}{2} \oplus c \cdot \omega)$
- ▶  $Apply(T, R_1) = \{0, \frac{1}{4}, \frac{1}{2}, \frac{3}{4}\}$
- ▶  $Apply(T, R_2) = \{\frac{1}{2}\}$

Therefore  $R_1 \not\sqsubseteq_{\text{pmay}} R_2$

Without internal choice in tests,  $R_1 \sqsubseteq_{\text{may}} R_2$



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- ▶  $T = a \cdot (\omega \cdot \frac{1}{2} \oplus \mathbf{0}) \sqcap (b \cdot \omega \cdot \frac{1}{2} \oplus c \cdot \omega)$
- ▶  $Apply(T, R_1) = \{0, \frac{1}{4}, \frac{1}{2}, \frac{3}{4}\} = \frac{1}{2} \cdot \{0, \frac{1}{2}\} + \frac{1}{2} \cdot \{0, 1\}$
- ▶  $Apply(T, R_2) = \{\frac{1}{2}\} = \frac{1}{2} \cdot \{\frac{1}{2}\} + \frac{1}{2} \cdot \{\frac{1}{2}\}$

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## Massacre of the axioms

$$a.(P \sqcap Q) \neq a.P \sqcap a.Q$$

$$P \neq P \square P$$

$$P \square (Q \sqcap R) \neq (P \square Q) \sqcap (P \square R)$$

$$P \sqcap (Q \square R) \neq (P \sqcap Q) \sqcap (P \sqcap R)$$

Hopes dashed:

$$P_{p\oplus} (Q \square R) \neq (P_{p\oplus} Q) \square (P_{p\oplus} R)$$

$$P \sqcap (Q_{p\oplus} R) \neq (P \sqcap Q)_{p\oplus} (P \sqcap R)$$

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All is not lost:

$$a.P \square a.Q = a.P \sqcap a.Q$$

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# Outline

Testing theory

Probabilistic CSP

The power of testing

**Simulations**

Equational theories

# Simulations: for may testing

Simulations relate states with distributions

## Definition (Segala)

- ▶  $\mathcal{R} \subseteq S_p \times \mathcal{D}(S_p)$  is a simulation if  
 $s \mathcal{R} \Delta$  and  $s \xrightarrow{\alpha} \Theta$  implies there exists some  $\Delta'$   
 $\Theta \overline{\mathcal{R}} \Delta'$  and  $\Delta \xrightarrow{\hat{\alpha}} \Delta'$

## Requirements:

- ▶ Lift  $S_p \mathcal{R} \mathcal{D}(S_p)$  to  $\mathcal{D}(S_p) \overline{\mathcal{R}} \mathcal{D}(S_p)$
- ▶ Generalise  $s \xrightarrow{\alpha} \Delta'$  to  $\Delta \xrightarrow{\hat{\alpha}} \Delta'$
- ▶  $\xrightarrow{\hat{\alpha}}$  means  $\xrightarrow{\tau}^* \xrightarrow{a} \xrightarrow{\tau}^*$
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## Lifting relation: $S \mathcal{R} D(S)$

$\Delta_1 \overline{\mathcal{R}} \Delta_2$  if

- ▶ source  $\Delta_1$  can be decomposed:

$$\Delta_1 = \sum_{i \in I} p_i \cdot \overline{s}_i$$

- ▶ each decomposed state  $s_i$  can be related:

$$s_i \mathcal{R} \phi_i, \text{ for some } \phi_i$$

- ▶ related  $\phi_i$  yield the target:

$$\Delta_2 = \sum_{i \in I} p_i \cdot \phi_i$$

## Example liftings: actions

$$(a.b \sqcap a.c)_{\frac{1}{2}} \oplus a.d \xrightarrow{a} b_{\frac{1}{2}} \oplus d$$

because  $(a.b \sqcap a.c) \xrightarrow{a} b$       $a.d \xrightarrow{a} d$

$$(a.b \sqcap a.c)_{\frac{1}{2}} \oplus a.d \xrightarrow{a} (b_{\frac{1}{2}} \oplus c)_{\frac{1}{2}} \oplus d$$

because

▶ source is  $\frac{1}{4} \cdot \overline{(a.b \sqcap a.c)} + \frac{1}{4} \cdot \overline{(a.b \sqcap a.c)} + \frac{1}{2} \cdot \overline{a.d}$

▶  $a.b \sqcap a.c \xrightarrow{a} \bar{b}$       $a.b \sqcap a.c \xrightarrow{a} \bar{c}$       $a.d \xrightarrow{a} \bar{d}$

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Lifting actions: 0 or 1 internal moves:  $\xrightarrow{\hat{\tau}}$

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## Results on simulations

$P \sqsubseteq_S Q$  - lifting of simulations to pCSP processes

- ▶  $\sqsubseteq_S$  is a precongruence for pCSP
  - ▶ For pCSP,  $P \sqsubseteq_S Q$  implies  $P \sqsubseteq_{\text{pmay}} Q$
  - ▶ For CSP,  $P \sqsubseteq_S Q$  iff  $P \sqsubseteq_{\text{pmay}} Q$
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- ▶  $P \sqsubseteq_S Q$  can be equationally characterised, over pCSP
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# Outline

Testing theory

Probabilistic CSP

The power of testing

Simulations

Equational theories

# Equations for CSP

tests contain probability choice, processes not

A selection:

$$\begin{aligned}
 a.P \sqcap a.Q &= a.P \sqcap a.Q \\
 a.P \sqcap (Q \sqcap R) &= (a.P \sqcap Q) \sqcap (a.P \sqcap R) \\
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For completeness add:  $P \sqsubseteq P \sqcap Q$      $P \sqcap Q = P \sqcap Q$     standard

Interesting derived equations:

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# Equations for pCSP

processes and tests contain probability choice

Standard:

$$\begin{aligned}
 P_p \oplus P &= P \\
 P_p \oplus Q &= Q_{1-p} \oplus P \\
 (P_p \oplus Q)_q \oplus R &= P_{p \cdot q} \oplus (Q_{\frac{(1-p) \cdot q}{1-p \cdot q}} \oplus R)
 \end{aligned}$$

Add:

$$\begin{aligned}
 a.(P_p \oplus Q) &\sqsubseteq a.P_p \oplus a.Q \\
 P \sqcap (Q_p \oplus R) &= (P \sqcap Q)_p \oplus (P \sqcap R)
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$$P_p \oplus Q = Q_{1-p} \oplus P$$

$$(P_p \oplus Q)_q \oplus R = P_{p \cdot q} \oplus (Q_{\frac{(1-p) \cdot q}{1-p \cdot q}} \oplus R)$$

Add:

$$a.(P_p \oplus Q) \sqsubseteq a.P_p \oplus a.Q$$

$$P \sqcap (Q_p \oplus R) = (P \sqcap Q)_p \oplus (P \sqcap R)$$

Interesting derived equations:

$$(P_p \oplus Q) \sqcap (P_p \oplus R) \sqsubseteq P_p \oplus (Q \sqcap R)$$

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# Equations for pCSP

processes and tests contain probability choice

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# The end

Lots of work to do:

- ▶ Prove relationship between simulations and may testing
- ▶ Must testing:
  - ▶ relate to *failure simulations*
  - ▶ equations ?
- ▶ Modal logics ?

## Related work

- ▶ Segala, Wang Yi, Jonsson, Larsen, Skou, Lowe, Cleaveland, Smolka, ...
- ▶ Google:
  - ▶ probabilistic process calculi: 631,000
  - ▶ testing probabilistic process calculi: 394,000
  - ▶ testing probabilistic processes: 4,800,000
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